

Aspects of 2-categorical logic

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“Isoregular theories, accessible 2-categories, and free constructions”

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Context, motivation, and goals

Accessible 1-categories

- ▶ Categories of models of (infinitary) first-order theories

See Makkai-Paré [1989]

Accessible 2-categories

- ▶ Lack and Rosický [2012]
- ▶ Lack and Tendas [2022, 2023]

Open problem: Show that 2-categories of interest are accessible with some limits

- ▶ Makkai [1997]
- ▶ Bourke [2021]: pseudo-algebra morphisms are key

Main contribution: A new method, many applications.

Plan

1. Review of 1-categorical logic
2. Isoregular theories
3. Main results and applications

Regular theories

$$\top, \quad s = t, \quad A \wedge B, \quad (\exists x: X)A(x).$$

Theorem. Let T be a regular theory.

- (i) **Syn**(T) is a regular category
- (ii) **Mod**(T) \cong **Reg**[**Syn**(T), **Set**]
- (iii) **Mod**(T) is accessible with finite products.

Examples.

- ▶ Theory of Abelian groups
- ▶ Theory of divisible Abelian groups: $(\exists y)ny = x$
- ▶ Theory of categories

Finite limit theories

$$\top, \quad s = t, \quad A \wedge B, \quad (\exists!x: X)A(x).$$

Theorem. Let \mathcal{T} be a finite limit theory.

- (i) **Syn**(\mathcal{T}) is a finite limit category
- (ii) **Mod**(\mathcal{T}) \cong **Lex**[**Syn**(\mathcal{T}), **Set**]
- (iii) **Mod**(\mathcal{T}) is finitely accessible with limits, i.e. locally finitely presentable.

Examples.

- ✓ Theory of Abelian groups
- ✗ Theory of divisible Abelian groups: $(\exists y)ny = x$
- ✓ Theory of categories

Isoregular 2-categories

Definition. A category \mathcal{C} is **regular** if it has finite limits and

1. every kernel pair has a coequaliser
2. regular epimorphisms are stable under pullback.

Definition. A 2-category \mathcal{K} is **isoregular** if it has finite 2-categorical limits and

1. every **fully faithful** kernel pair has a coequaliser
2. **fully faithful** regular epimorphisms are stable under pullback.

A kernel pair

$$\Delta_f \begin{array}{c} \xrightarrow{\pi_1} \\ \xrightarrow{\pi_2} \end{array} A \xrightarrow{f} B$$

is said to be **fully faithful** if π_1 (or π_2) is fully faithful.

Characterisation

Proposition. Let $f: A \rightarrow B$ be a morphism in \mathcal{K} . The following are equivalent:

- (i) the morphism f is an FFK-morphism, i.e. it has a fully faithful kernel pair
- (ii) for every $X \in \mathcal{K}$, $\mathcal{K}(X, f): \mathcal{K}(X, A) \rightarrow \mathcal{K}(X, B)$ has a fully faithful kernel pair

(iii) the diagram
$$\begin{array}{ccc} A & \xlongequal{\quad} & A \\ \parallel & & \downarrow f \\ A & \xrightarrow{f} & B \end{array}$$
 is a semistrict pullback

- (iv) the morphism f is faithful and full on identities
- (v) the unique $r: A \rightarrow \Delta_f$ such that $\pi_1 \circ r = 1_A$ and $\pi_2 \circ r = 1_A$ is an equivalence
- (vi) the morphism $f: A \rightarrow B$ is a subcontractible object of the slice 2-category $\mathcal{K}_{/B}$

Example. $\{ (d, c, \lambda: \Delta c \Rightarrow d) \mid d \in C', c \in C, (c, \lambda) \text{ limit of } d \} \xrightarrow{\pi_1} C'.$

Factorisation

Proposition. Let \mathcal{K} be an isoregular 2-category. Every FFK-morphism factors as a fully faithful regular epimorphism followed by a monomorphism.

Idea: we can form

$$\begin{array}{ccccc} \{x: X, y: Y \mid B(x, y)\} & \xrightarrow{\quad} & X \times Y & \xrightarrow{\pi_1} & X \\ & \searrow & & \nearrow & \\ & & \{x: X \mid (\exists y: Y) B(x, y)\} & & \end{array}$$

if the fibers $\{y: Y \mid B(x, y)\}$, for $x: X$, are subcontractible (cf. HoTT).

Example: $\{d \in C^I \mid \exists (c, \lambda) \text{ limit of } d \text{ in } C\} \twoheadrightarrow C^I$.

Isoregular theories

Enriched categorical logic for $\mathcal{V} = \mathbf{Cat}$. [Rosický–Tendas 2025 & 2026]

Arities = finitely presented categories.

Definition. A **finitary language** \mathbb{L} is given by:

- ▶ basic sorts: S, T, \dots
- ▶ function symbols: $f: S_1^{\mathbb{C}_1}, \dots, S_n^{\mathbb{C}_n} \rightarrow S^{\mathbb{C}}$
- ▶ relation symbols: $R \multimap S_1^{\mathbb{C}_1}, \dots, S_n^{\mathbb{C}_n}$

Example. The language $\mathbb{L}_{\mathbf{GFib}}$ for Grothendieck fibrations:

- ▶ two basic sorts: E and B
- ▶ one function symbols: $p: E \rightarrow B$
- ▶ one relation symbol: $\text{Cart} \multimap E^{[1]}$

Isoregular theories

Forms of judgement: $s : X$, $A : \text{prop}$, $A_1, \dots, A_n \vdash A$

Propositions: \top , $s = t$, $A \wedge B$, $(\exists x : X)A(x)$, $A^{[1]}$

Deduction rules: for example

$$\frac{(x : X, y : Y) \quad B(x, y) : \text{prop} \quad \mathcal{J}_{\text{faithful}}(B) \quad \mathcal{J}_{\text{full-id}}(B)}{(x : X) \quad (\exists y : Y)B(x, y) : \text{prop},}$$

$$\frac{A_1, \dots, A_n \vdash A}{A_1^{[1]}, \dots, A_n^{[1]} \vdash A^{[1]}}$$

Example: The theory $\mathbb{T}_{\mathbf{GFib}}$ has

$$(y : E, u : B^{[1]}) \quad \text{cod}(u) = p(y) \vdash (\exists v : E^{[1]}) \text{Cart}(v) \wedge \text{cod}(v) = y \wedge p^{[1]}(v) = u$$

Main results

Theorem. Let T be an isoregular theory.

- (i) **Syn**(T) is an isoregular 2-category
- (ii) **Mod**(T) \cong **IsoReg**[**Syn**(T), **Cat**]
- (iii) **Mod**(T) is accessible with flexible limits.

Proof.

- (i) Use deduction rules for isoregular theories.
- (ii) Direct calculation.
- (iii) 2-dimensional category theory. □

Applications: general category theory

Theorem.

- (i) The 2-categories **Lex**, **Rex**, **RLex** are accessible with flexible limits.
- (ii) The forgetful 2-functors

$$\begin{array}{ccc} \mathbf{RLex} & \longrightarrow & \mathbf{Lex} \\ \downarrow & & \downarrow \\ \mathbf{Rex} & \longrightarrow & \mathbf{Cat} \end{array}$$

have left biadjoints (they are accessible and preserve flexible limits).

Proof of (i). Write $\mathbb{T}_{\mathbf{Lex}}, \dots$

- ▶ Finite case of Joyal's bicompletion

Applications: categorical logic

Theorem.

- (i) The 2-categories **GFib**, **Cmp**, **Clan** are accessible with flexible limits.
- (ii) The forgetful 2-functors

$$\mathbf{Cmp} \rightarrow \mathbf{GFib}, \quad \mathbf{Clan} \rightarrow \mathbf{Cat}$$

have left biadjoints (they are accessible and preserve flexible limits).

Proof of (i). Write $\mathbb{T}_{\mathbf{GFib}}, \dots$

Note. Many other applications:

- ▶ multicategories (**Mult**, **RepMult**, ...).
- ▶ categorical algebra (**Reg**, **Ex**, **Mal**, **SmAb**, ...)

Summary

- ▶ Isoregular 2-categories & isoregular theories
- ▶ 2-categories of models are accessible with flexible limits
- ▶ Applications to category theory, categorical logic, categorical algebra

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