

Monoidal bicategories, differential linear logic , and analytic functors *

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What do logicians have in common? *

Theorem (Completeness for propositional logic)

$\Gamma \vdash A$

$\S$

defined syntactically
via calculus

- Hilbert system
- Natural deduction
- Sequent calculus

iff

$\Gamma \models A$

$\S$

defined semantically
via truth values

* Alex Wilkie, Lc2010.

Deduction rules

$$\frac{}{A \vdash A} \text{ax}$$

$$\frac{\Gamma \vdash A \quad \Delta \vdash B}{\Gamma, \Delta \vdash A \wedge B} \wedge_R$$

$$\frac{\Gamma, A, B \vdash C}{\Gamma, A \wedge B \vdash C} \wedge_L$$

Exercise Prove $A \vdash A \wedge A$

Solution

$$\frac{\frac{\frac{}{A \vdash A}}{A, A \vdash A \wedge A} \quad \frac{}{A \vdash A}}{A \vdash A \wedge A} \text{c}$$

Exercise Prove $A \wedge B \vdash A$

Solution

$$\frac{\frac{}{A \vdash A}}{A, B \vdash A} \text{w}$$
$$\frac{}{A \wedge B \vdash A}$$

Linear Logic

- Drop structural rules for arbitrary formulas.

$$\frac{}{A \vdash A} ax \qquad \frac{\Gamma \vdash A \quad \Delta \vdash B}{\Gamma, \Delta \vdash A \otimes B} \otimes_R \qquad \frac{\Gamma, A, B \vdash C}{\Gamma, A \otimes B \vdash C} \otimes_L$$

- Note $A \not\vdash A \otimes A$, $A \otimes B \vdash A$

- We think of hypotheses as resources

- This is the difference between

$(\mathbb{L}, \times, 1)$
cartesian, e.g. Set

$(\mathbb{L}, \otimes, I)$
monoidal, e.g. Vect_k

Proofs - as - maps

Linear Logic

$$\frac{}{A \vdash A} ax$$

$$\frac{\Gamma \vdash A \quad \Delta \vdash B}{\Gamma, \Delta \vdash A \otimes B} \otimes_R$$

$$\frac{\Gamma \vdash A \quad \Delta, A \vdash B}{\Gamma, \Delta \vdash B} cut$$

Monoidal category

$$A \xrightarrow{id_A} A$$

$$\frac{\Gamma \xrightarrow{a} A \quad \Delta \xrightarrow{b} B}{\Gamma \otimes \Delta \xrightarrow{a \otimes b} A \otimes B}$$

$$\frac{\Gamma \xrightarrow{a} A \quad \Delta \otimes A \xrightarrow{b} B}{\Gamma \otimes \Delta \xrightarrow{a \otimes 1} A \otimes \Delta \xrightarrow{b'} B}$$

Commutative diagrams as proof equations

Functoriality of $\otimes$ corresponds to:

$$\frac{\frac{\Gamma \vdash A \quad \Delta \vdash B}{\Gamma, \Delta \vdash A \otimes B} \quad \frac{A \vdash A' \quad B \vdash B'}{A \otimes B \vdash A' \otimes B'}}{\Gamma, \Delta \vdash A' \otimes B'} \text{ cut}$$



$$\frac{\frac{\Gamma \vdash A \quad A \vdash A'}{\Gamma \vdash A'} \text{ cut} \quad \frac{\Delta \vdash B \quad B \vdash B'}{\Delta \vdash B'} \text{ cut}}{\Gamma, \Delta \vdash A' \otimes B'}$$

Linear Logic

Idea : structural rules for some formulas, via **modality** :

"of course A"
"bang A"

!A

we can use !A as
many times as we want

$$\frac{\Gamma, A \vdash B}{\Gamma, !A \vdash B} \text{d}$$

$$\frac{! \Gamma \vdash A}{! \Gamma \vdash !A} \text{p}$$

$$\frac{\Gamma \vdash B}{\Gamma, !A \vdash B} \text{w}$$

$$\frac{\Gamma, !A, !A \vdash B}{\Gamma, !A \vdash B} \text{c}$$

Note

!A $\vdash$ **!A** $\otimes$ **!A** ,

!A $\otimes$ **!B** $\vdash$ **!A**

Linear exponential comonad

$(\mathbb{L}, \otimes, I)$ a symmetric monoidal category with

$$\bullet \mathbb{L} \xrightarrow{!} \mathbb{L}$$

$$\bullet \text{ maps } d_A : !A \rightarrow A$$

$$\bullet \text{ maps } w_A : !A \rightarrow I$$

$$\bullet \text{ maps } p_A : !A \rightarrow !!A$$

$$\bullet \text{ maps } c_A : !A \rightarrow !A \otimes !A$$

KEY

IDEA

maps $f: A \rightarrow B$ in $\mathbb{L}$ are linear

maps $f: !A \rightarrow B$ in $\mathbb{L}$ are

non-linear

Differential Linear Logic (T. Ehrhard & L. Reiguiet)

Maps in many models of Linear Logic admit
a well-behaved differentiation operation

⇒ introduce a syntax for it

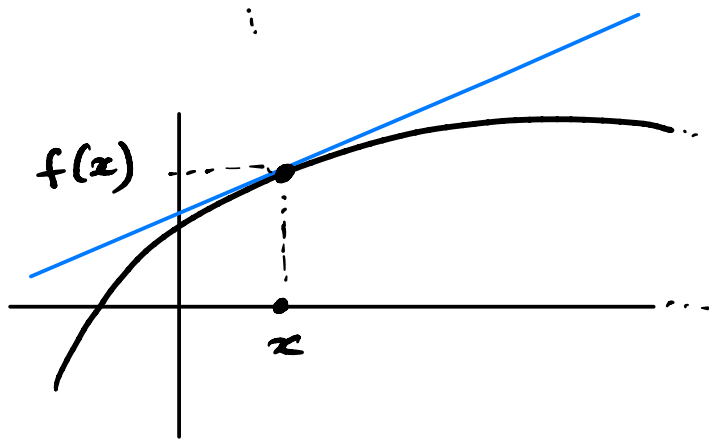
⇒ differential λ -calculus, differential linear logic

⇒ applications to ordinary λ -calculus
(via Taylor series)

2. Differential Linear Logic

Recall For $f: \mathbb{R}^m \longrightarrow \mathbb{R}^n$, $\vec{x} \in \mathbb{R}^m$,

the Jacobian $Df_{\vec{x}}$ is an $(m \times n)$ -matrix



Then $f'(x)e_j$ is the j th column vector of $[f'(x)]$, and (27) shows therefore that the number $(D_j f_i)(x)$ occupies the spot in the i th row and j th column of $[f'(x)]$. Thus

$$[f'(x)] = \begin{bmatrix} (D_1 f_1)(x) & \cdots & (D_n f_1)(x) \\ \vdots & \ddots & \vdots \\ (D_1 f_m)(x) & \cdots & (D_n f_m)(x) \end{bmatrix}.$$

If $\mathbf{h} = \sum h_j \mathbf{e}_j$ is any vector in $\mathbb{R}^n$, then (27) implies that

$$(30) \quad f'(x)\mathbf{h} = \sum_{i=1}^m \left\{ \sum_{j=1}^n (D_j f_i)(x) h_j \right\} \mathbf{u}_i.$$

$\Rightarrow$

$Df_{\vec{x}}(\vec{a})$

is

linear in $\vec{a} \in \mathbb{R}^m$

non-linear in $\vec{x} \in \mathbb{R}^m$

Differential categories (Blute, Cockett, Seely)

Let $(\mathbb{L}, \otimes, \mathbb{I})$ with $(!, p, d)$ as before.

WANT An operator $D[-]$

$$f : !A \longrightarrow B$$

to

$$D[f] : A \otimes !A \longrightarrow B$$

linear non-linear

subject to usual axioms of differential calculus:

constant, sum, product, chain.

Theorem (Fiore, Blute & Cochet & Lemay & Seely)

A differential category is determined by maps

$$\bar{d}_A : A \longrightarrow !A \quad \text{satisfying}$$

$$\begin{array}{ccc} A & \xrightarrow{\bar{d}} & !A \\ & \searrow 1_A & \downarrow d \\ & & A \end{array}$$

$$\begin{array}{ccccc} A \otimes I & \xrightarrow{\cong} & A & \xrightarrow{d} & !A \\ \bar{d} \otimes \bar{w} \downarrow & & & & \downarrow p \\ !A \otimes !A & \xrightarrow{\bar{d} \otimes p} & !!A \otimes !!A & \xrightarrow{\bar{c}} & !!A \end{array}$$

$$\begin{array}{ccc} A \otimes !B & \xrightarrow{\bar{d} \otimes 1} & !A \otimes !B \\ 1 \otimes d \downarrow & & \downarrow w^2 \\ A \otimes B & \xrightarrow{\bar{d}} & !(A \otimes B) \end{array}$$

Note

For $f : !A \rightarrow B$, get $D[f] : A \otimes !A \rightarrow B$ as

$$A \otimes !A \xrightarrow{\bar{d} \otimes 1} !A \otimes !A \xrightarrow{\bar{c}} !A \xrightarrow{f} B$$

3. Analytic functors

Calculus

$$\mathbb{R} \xrightarrow{f} \mathbb{R} \quad \text{linear}$$

$$\mathbb{R}^m \longrightarrow \mathbb{R}^n \quad \text{linear}$$

$$\mathbb{R} \xrightarrow{f} \mathbb{R}$$

$$\mathbb{R}^m \longrightarrow \mathbb{R}^n$$

Functors

$$\underline{\text{Set}} \xrightarrow{F} \underline{\text{Set}} \quad \text{linear functor}$$

$$\underline{\text{Set}}^{A^{\text{op}}} \longrightarrow \underline{\text{Set}}^{B^{\text{op}}} \quad \text{linear functor}$$

$$\underline{\text{Set}} \xrightarrow{F} \underline{\text{Set}} \quad \text{analytic}$$

(Joyal '86)

$$\underline{\text{Set}}^{A^{\text{op}}} \longrightarrow \underline{\text{Set}}^{B^{\text{op}}} \quad \text{analytic}$$

(FGHW '08)

Linear maps and Linear Functors

Definition Let A, B be small categories. A

linear map^{*} $M : A \rightarrow B$ is a functor

$$\begin{array}{ccc} B^{\text{op}} \times A & \xrightarrow{M} & \underline{\text{Set}} \\ (b, a) & \longmapsto & M(b, a) \end{array}$$

the (b, a) -entry of the matrix

The linear functor $\tilde{M} : \underline{\text{Set}}^{A^{\text{op}}} \rightarrow \underline{\text{Set}}^{B^{\text{op}}}$

$$(\tilde{M}x)_b = \sum_{a \in A} M(b, a) \times x_a$$

$\nearrow \cong$
a quotient

* Bénabou's profunctors.

Towards Linear Logic

$\Downarrow$
Lin = the category of

- small categories
- matrices

Prop The category Lin admits a symmetric monoidal structure where $A \otimes B =_{\text{def}} A \times B$.

(It has also biproducts and is compact closed)

The linear exponential comonad

For $A \in \underline{\text{Lin}}$, define $!A \in \underline{\text{Lin}}$ as the free symmetric monoidal category on A :

- **objects**: $\vec{a} = (a_1, \dots, a_n)$, for $n \in \mathbb{N}$.
- **maps**: $(\sigma, \vec{f}) : (a_1, \dots, a_n) \rightarrow (a'_1, \dots, a'_n)$
where $\sigma \in \Sigma_n$, $f_i : a_i \rightarrow a'_{\sigma(i)}$

Theorem [FGH] This definition extends to a linear exponential comonad on $\underline{\text{Lin}}$.

Analytic functors

Definition Let A, B be small categories. A

non-linear map $F : A \longrightarrow B$ is a linear map

$F : !A \longrightarrow B$, ie. a functor $F : B \times !A^{\varphi} \longrightarrow \underline{\text{Set}}$

The analytic functor $\tilde{F} : \underline{\text{Set}}^{A^{\varphi}} \longrightarrow \underline{\text{Set}}^{B^{\varphi}}$ is

$$(\tilde{F} \times)_b = \sum_{\vec{a} \in !A} F[b; \vec{a}] \times \prod_{i=1}^n X^{a_i} \Big/ \sim$$

"coefficients"
=_{\text{def}} X(a_1) \times \dots \times X(a_n)
quotient

Joyal's analytic functors

Let $A = B = 1$. Then $!1 = \mathbb{P} =$ category of

- natural numbers
- permutations

$$F : !1 \rightarrow 1 = F : \mathbb{P} \rightarrow \underline{\text{Set}} \quad \text{and}$$

$$\tilde{F} : \underline{\text{Set}} \longrightarrow \underline{\text{Set}}$$

$$X \longmapsto \sum_{n \in \mathbb{N}} F[n] \times X^n / \simeq_{\sigma_n}$$

Note Cf. $f(x) = \sum_{n \in \mathbb{N}} f_n \frac{x^n}{n!}$.

Differentiation

- For $f(x) = \sum_n f_n \frac{x^n}{n!}$, $f'(x) = \sum_n f_{n+1} \frac{x^n}{n!}$

- For $F: \underline{\text{Set}} \longrightarrow \underline{\text{Set}}$

$$F'(X) = \sum_{n \in \mathbb{N}} F[n+1] \times X^n / \simeq_{\sigma_n} \quad (\text{Joyal})$$

- For $F: \underline{\text{Set}}^{A^{\text{op}}} \longrightarrow \underline{\text{Set}}^{B^{\text{op}}}$

What should we do ?

Differentiation

Define $\bar{d}_A : A \dashrightarrow !A$ by

$$\begin{aligned} !A \times A^\omega &\xrightarrow{\bar{d}_A} \underline{\text{Set}} \\ (\vec{a}, a) &\longmapsto !A[\vec{a}, (a)] \end{aligned}$$

Theorem [FGH] This satisfies the axioms
for a differential category.

Differentiation

For $F: \underline{\text{Set}}^{A^{\text{op}}} \longrightarrow \underline{\text{Set}}^{B^{\text{op}}}$ analytic, $X \in \underline{\text{Set}}^{A^{\text{op}}}$

the Jacobian $dF(X): A \dashrightarrow B$ is given by

$$dF(X)(b, a) = \sum_{\vec{a} \in !A} F[b; \underbrace{\vec{a} \oplus (a)}_{\substack{\text{length } n+1 \\ \text{if } \vec{a} \text{ has length } n}}] \times X^{\vec{a}}$$

Corollary All the rules of calculus hold
for analytic functors.

Monoidal bicategories

The example of analytic functors does not fit into the theory of monoidal categories since it is inherently 2-categorical, e.g.

$$\begin{array}{ccc} & ((A \otimes B) \otimes C) \otimes D & \\ \swarrow a & & \searrow a \\ (A \otimes (B \otimes C)) \otimes D & \Downarrow \pi & (A \otimes B) \otimes (C \otimes D) \\ \searrow a & & \swarrow a \\ A \otimes ((B \otimes C) \otimes D) & \xrightarrow{a} & (A \otimes (B \otimes (C \otimes D))) \end{array}$$

$\Rightarrow$ The paper develops the theory of monoidal bicategories (Kapranov-Voevodsky, ...)

Coherence and rewriting

$$\begin{array}{c}
 \frac{!A \xrightarrow{f} B}{!A \xrightarrow{f^*} !B} \quad \frac{!B \xrightarrow{d_B} B}{!B \xrightarrow{d_B^*} !B} \\
 \hline
 !A \xrightarrow{d_B^* f^*} !B \quad \text{cut}
 \end{array}$$



$$\begin{array}{c}
 \frac{!A \xrightarrow{f} B}{!A \xrightarrow{f^*} !B} \quad !B \xrightarrow{d_B} B \\
 \hline
 !A \xrightarrow{d_B f^*} B \quad \text{cut} \\
 \hline
 !A \xrightarrow{(d_B f^*)^*} !B
 \end{array}$$



$$\frac{!A \xrightarrow{f} B}{!A \xrightarrow{f^*} !B}$$

$$\begin{array}{c}
 \frac{!A \xrightarrow{f} B}{!A \xrightarrow{f^*} !B} \quad !B \xrightarrow{1} !B \\
 \hline
 !A \xrightarrow{f^*} !B \quad \text{cut}
 \end{array}$$



Summary

- The bicategory Lin of small categories and linear maps can be equipped with the structure of a model of differential linear logic.
- This gives an extension of Joyal's differential calculus for analytic functors to 'many variables'.