

# Prefix monoids of groups and right units of special inverse monoids

Igor Dolinka

*Department of Mathematics and Informatics, University of Novi Sad, Serbia*

36th NBSAN Meeting, a satellite event to the 75th BMC

Manchester, UK, 20 June 2024



## Joint work with Robert D. Gray (UEA)

I.Dolinka, R.D.Gray: *Prefix monoids of groups and right units of special inverse monoids*, Forum of Mathematics, Sigma **11** (2023), Article e97, 19 pp.

# The “driving engine” (part I)

H. H. Wilhelm Magnus (1930/31):

*The word problem for every one-relator group  $G_p\langle A \mid r = 1 \rangle$  is decidable.*

**Reason** (the Magnus-Moldavansky hierarchy):

- ▶  $G = G_p\langle A \mid r = 1 \rangle$  embeds into an HNN-extension of its (f.g.) subgroup  $L = G_p\langle A' \mid r' = 1 \rangle$  w.r.t. a pair of free (“Magnus”) subgroups of  $L$ , where  $|r'| < |r|$ ;
- ▶ This suffices to reduce the WP for  $G$  to that of  $L$ ;
- ▶ Eventually, we end up with a free group of finite rank, where we trivially solve the WP.
- ▶ NB. There is an older approach (Magnus’ original) using amalgamated free products.

## The “driving engine” (part II)

**Open problem** (as of 20 June 2024):

*Does every one-relator monoid  
 $\text{Mon}\langle A \mid u = v \rangle$  have a decidable WP?*

**S.I.Adian** (1966) – The word problem for  $\text{Mon}\langle A \mid u = v \rangle$  is decidable for:

- ▶ *special monoids* – the def. relation is of the form  $u = 1$ ,
- ▶ the case when both  $u, v$  are *non-empty*, and have different *initial* letters and different *terminal* letters.

**Adian & Oganessian** (1987) – The general problem reduces to two particular cases:

- ▶  $\text{Mon}\langle a, b \mid aUb = aVa \rangle$ ,
- ▶  $\text{Mon}\langle a, b \mid aUb = a \rangle$  (the “monadic” case).

NB. These presentations define **right cancellative** monoids.

## The Lead Role #1: Prefix monoids (in groups)

Let  $G = \text{Gp}\langle A \mid w_i = 1 \ (i \in I) \rangle$  be a group.

The **prefix monoid** of this group (presentation) = the submonoid of  $G$  generated by the elements represented by all prefixes of all  $w_i$ 's

The prefix monoid is **dependent on the concrete presentation** of  $G$  – one fixed (isomorphism type of a) group can have many presentations, leading to many prefix monoids.

**Prefix Membership Problem** (PMP): Given a word over  $A \cup A^{-1}$ , decide whether it represents an element of the prefix monoid (w.r.t. the given group presentation)

**IgD & RDG** (TrAMS, 2021): A kaleidoscope of sufficient conditions (via amalgamated products and HNN extensions) ensuring **decidability** for the PMP

## The Lead Role #2: Right units (in inverse monoids)

Let  $M$  be an **inverse** monoid.

$r \in M$  is a **right unit**  $\iff r \mathcal{R} 1 \iff rr^{-1} = 1$

Fun facts:

- ▶ Right units of  $M$  form a **right cancellative submonoid**  $R$  of  $M$ .
- ▶ If  $M = \text{Inv}\langle A \mid w_i = 1 \ (i \in I) \rangle$  (i.e.  $M$  is a **special** inverse monoid) then  $R$  is generated by elements represented by all prefixes of all  $w_i$ 's.
- ▶ So, in the natural map  $M \rightarrow G = \text{Gp}\langle A \mid w_i = 1 \ (i \in I) \rangle$ , the RU-monoid  $R$  of  $M$  is mapped **onto** the prefix monoid of  $G$ .
- ▶ If  $M$  happens to be  $E$ -unitary, the restriction of this map to  $R$  is a monoid **isomorphism**.
- ▶ Consequently, the RU-monoid of any  $E$ -unitary special inverse monoid (SIM) is **group-embeddable**.

## The “driving engine” (part III)

Ivanov, Margolis & Meakin (JPAA, 2001):

The (right cancellative) monoid  $\text{Mon}\langle A \mid aUb = aVc \rangle$  ( $b \neq c$ ) embeds (as the monoid of right units) into

$$\text{Inv}\langle A \mid aUbc^{-1}V^{-1}a^{-1} = 1 \rangle.$$

Similarly,  $\text{Mon}\langle A \mid aUb = a \rangle$  embeds into  $\text{Inv}\langle A \mid aUba^{-1} = 1 \rangle$ .

Hence, the WP for one-relator monoids reduces to the WP for **one-relator inverse monoids**.

Fun facts: when  $w$  is **cyclically reduced** then

- ▶  $\text{Inv}\langle A \mid w = 1 \rangle$  is  $E$ -unitary;
- ▶ the WP for  $\text{Inv}\langle A \mid w = 1 \rangle$  reduces to the PMP for  $\text{Gp}\langle A \mid w = 1 \rangle$ .

# Surprise, surprise...!

RDG (Inventiones, 2020):

*There **exists** a one-relator special inverse monoid  
with an undecidable WP. [!!!]*

Fun facts:

- ▶ the counterexample(s) is/are even **E-unitary**;
- ▶ at the heart of the proof is **Lohrey-Steinberg's** result (JAlg, 2008) that the **RAAG**  $A(P_4)$  has a fixed f.g. submonoid with undecidable membership;
- ▶ then,  $A(P_4)$  **embeds** into a one-relator group  $G = \text{Gp}\langle a, b \mid \dots \rangle$ ;
- ▶ finally, a one-relator SIM  $M = \text{Inv}\langle a, b, t \mid \dots \rangle$  is constructed so that  $u \in \{a, b, a^{-1}, b^{-1}\}^*$  represents an element of the “critical” undecidable f.g. submonoid of  $G \iff tut^{-1}$  is a **right unit** in  $M$ .

Still, this **does not** invalidate the IMM approach.



# Know your limits

Guba (1997):

For any monadic  $M = \text{Mon}\langle a, b \mid aUb = a \rangle$  constructs  $G_M = \text{Gp}\langle x, y, A \mid xWyx^{-1} = 1 \rangle$  (for some  $W \in (A \cup \{x, y\})^*$ ) such that the WP for  $M$  reduces to PMP for  $G_M$ .

However, there are groups  $G = \text{Gp}\langle A \mid w = 1 \rangle$  with:

- ▶  $w$  reduced and undecidable PMP for  $G$  (IgD, RDG, 2021);
- ▶  $w = uv^{-1}$  reduced ( $u, v \in A^+$ ) and undecidable PMP for  $G$  (Foniqi, RDG, CFNB, to appear);
- ▶  $w \in A^+$  and undecidable submonoid membership problem for  $G$  (again, FGNB).

## Mon vs Inv

Obviously (imagine **Snape**'s voice here), one-relator/f.p. special monoids and special inverse monoids are very different creatures. For example:

- ▶ the group of units  $U$  of a  $M = \text{Mon}\langle A \mid w = 1 \rangle$  is a one-relator/f.p. group;
- ▶ the RU-monoid of  $M$  is a free product of  $U$  and a free monoid of finite rank;
- ▶ all other maximal subgroups of  $M$  are  $\cong U$ .

In contrast:

- ▶ the group of units  $U$  of a  $M = \text{Inv}\langle A \mid w = 1 \rangle$  can be non-one-relator (**RGD**, **Ruškuc**, Jussieu, to appear);
- ▶ the RU-monoid of  $M$  can be even non-f.p.;
- ▶ other maximal subgroups of  $M$  can be wildly different from  $U$ .

# The questions

All of this very much justifies the study of **prefix monoids in f.p. groups** and **RU-monoids in f.p. SIMs** in their own right.

- (1) What can the prefix monoids of f.p. groups be?
- (2) What can the RU-monoids of f.p. SIMs be?
- (3) What are the possible groups of units of these monoids?
- (4) What are the possible Schützenberger groups of these monoids?

## Recursive stuff

A group  $G$  is **recursively presented** if

$$G = \text{Gp}\langle A \mid w_i = 1 \ (i \in I) \rangle$$

where  $A$  is finite and  $\{w_i : i \in I\}$  is a **r.e. language** over  $A \cup A^{-1}$ .

Similarly, a **monoid** is recursively presented if

$$M = \text{Mon}\langle A \mid u_i = v_i \ (i \in I) \rangle$$

where  $A$  is finite and  $\{(u_i, v_i) : i \in I\}$  is a **r.e. subset** of  $A^* \times A^*$ .

**The Higman Embedding Theorem:** A finitely generated group embeds into a f.p. group if and only if it is recursively presented.

- 👉 A finitely generated monoid embeds into a f.p. group if and only if it is group-embeddable and recursively presented.
- 👉 Every prefix monoid (of a f.p. group) is f.g.  
 $\implies$  it is recursively presented.

# The characterisation of prefix monoids (of f.p. groups)

Two (easy) facts:

- ▶ Every group-embeddable **f.p.** monoid arises as a prefix monoid.
- ▶ If a **group** arises as a prefix monoid then it is **f.p.** So, not all group-embeddable recursively presented monoids are prefix monoids.

**Theorem** (IgD, RDG, 2023):

*For every group-embeddable recursively presented monoid  $M$  there is a natural number  $\mu_M$  such that*

$$M * \Sigma_k^*$$

*is a prefix monoid (with  $|\Sigma_k| = k$ ) if and only if  $k \geq \mu_M$ .*

**Also:**

*The class of groups of units of prefix monoids is precisely the recursively presented groups.*

## Recursively enumerable stuff

Let  $G$  be a f.p. group (generated by  $A$ ). Let  $L \subseteq (A \cup A^{-1})^*$  be a **recursively enumerable language** such that the set of all elements of  $G$  represented by words from  $L$  forms a subgroup  $H \leq G$ . Then  $H$  is said to be a **recursively enumerable subgroup** of  $G$ .

**NB.** A r.e. subgroup of  $G$  is not necessarily finitely generated. However, all f.g. (i.e. recursively presented) subgroups of  $G$  are r.e.

**Theorem** (IgD, RDG, 2023):

*A group  $H$  arises as a Schützenberger group of a prefix monoid (of a f.p. group)  $\iff H$  arises as a r.e. subgroup of a f.p. group.*

Ingredients:

- ▶  $M$  (left/right) cancellative  $\implies$  every Sch-group **embeds** into the group of units of  $M$ .
- ▶ For every r.e. subgroup  $H$  of a f.p. group  $G$  there is a f.p. overgroup  $G_1 \geq G$  and  $t \in G_1$  such that  $G \cap t^{-1}Gt = H$ .

# RU-monoids (take 1)

Again, some (easy) facts:

- ▶ Every RU-monoid is a **right cancellative** recursively presented monoid.
- ▶ If the monoid of right units of a f.p. SIM is a group  $\implies$  it is f.p.

**Theorem 1** (RDG, Kambites, JEMS, to appear):

*The class of groups of units of f.p. SIMs (and thus of RU-monoids) is precisely the recursively presented groups.*

**Theorem 2** (RDG, Kambites):

*A group arises as a maximal subgroup (i.e. as a group  $\mathcal{H}$ -class) of a f.p. SIM  $\iff$  it arises as a r.e. subgroup of a f.p. group.*

# RC-presentations

$$M = \text{MonRC}\langle A \mid \mathfrak{R} \rangle$$

$\Leftrightarrow M \cong A^*/\mathfrak{R}^{\text{RC}}$ , where  $\mathfrak{R}^{\text{RC}}$  is the intersection of all congruences  $\sigma$  of  $A^*$  such that

- ▶  $\mathfrak{R} \subseteq \sigma$ ,
- ▶  $A^*/\sigma$  is right cancellative.

**A.J.Cain** (2005) (+ Robertson, Ruškuc, 2008): A concept of formal, syntactic derivation for RC-presentations.

**Theorem** (IgD, RDG, 2023):

*Every finitely RC-presented monoid is an RU-monoid.*

In a way, this is a generalisation of the Ivanov-Margolis-Meakin result.



## RU-monoids (take 2)

**Theorem** (IgD, RDG, 2023):

*The class of Schützenberger groups of RU-monoids is exactly the class of r.e. subgroups of f.p. groups.*

**Open Problem:** Characterise the class of all RU-monoids.

In the remainder of the talk, I'll present **two interesting phenomena** in this vein discovered by IgD+RDG during this Spring's online sessions.

# The RU-monoid in the Gray-Ruškuć construction (1)

**RDG, Ruškuć:** For any group  $G$  (f.p. or not) and f.g. submonoid  $T$  of  $G$ , a(n  $E$ -unitary) SIM  $M$  is constructed (which is f.p. when  $G$  is) such that:

- ▶  $U(M) \cong G * U(T)$ ,
- ▶ if the monoid of right units of  $M$  is f.p. so must be both  $G$  and  $T$ .

With the right choice of parameters, this produces:

- ▶ a one-relator SIM whose group of units is **not one-relator**;
- ▶ a one-relator SIM whose group of units is f.p. but whose RU-monoid is **not f.p.**;
- ▶ a f.p. SIM whose group of units is **not f.p.**

## The RU-monoid in the Gray-Ruškcuc construction (2)

IgD, RDG (2024):

The RU-monoid of  $M$  = the greatest right cancellative image of the HNN-like Otto-Pride extension of  $G$  w.r.t.  $T \hookrightarrow G =$

$$\text{MonRC}\langle A, B, t \mid u_i = v_i (i \in I), tw_j = b_j t (j \in J) \rangle$$

where  $G = \text{Mon}\langle A \mid u_i = v_i (i \in I) \rangle$  and  $T = \langle w_j : j \in J \rangle_G$ .

Hence:

- ▶ If  $G$  is f.p. then the RU-monoid of  $M$  is necessarily **finitely RC-presented**;
- ▶ The group of units  $U(M)$  can still be **not f.p.**, and also the RU-monoid can be **not f.p.** (as a monoid!);
- ▶ There is a **finitely RC-presented** monoid  $S$  in which the complement of the group of units  $S \setminus U$  is an **ideal**, and still  $U$  is **not f.p.**

**Conclusion:** RC-presentations are strange animals!

# The Gray-Kambites construction (1)

Realising an arbitrary recursively presented group as the **group of units** of a f.p. SIM. Here we present a slight generalisation (by IgD & RDG).

$$T = \text{MonRC}\langle A \mid u_i = v_i \ (i = 1, \dots, k) \rangle$$

$S = \langle B \rangle_T$  – a f.g. submonoid

$M_{T,S}$  – a f.p. SIM gen. by  $A$  and  $p_0, p_1, \dots, p_k, z, d$  subject to

$$p_i a p_i^{-1} p_i a^{-1} p_i^{-1} = 1 \quad (a \in A, i = 0, 1, \dots, k)$$

$$p_i u_i d^{-1} v_i^{-1} p_i^{-1} = 1 \quad (i = 1, \dots, k)$$

$$p_0 d p_0^{-1} = 1$$

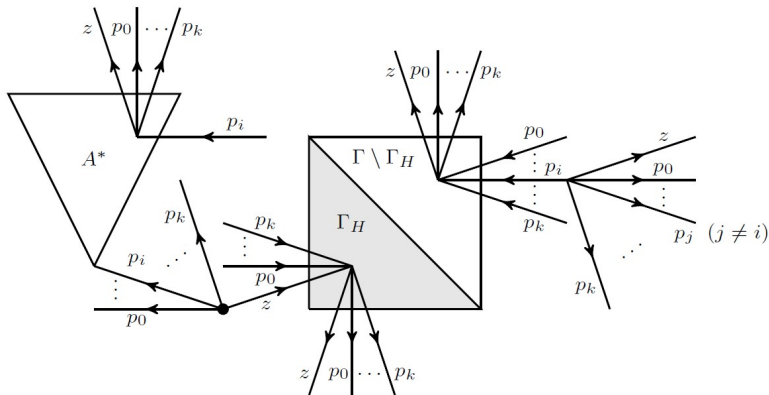
$$z b z^{-1} z b^{-1} z^{-1} = 1 \quad (b \in B)$$

$$z \left( \prod_{i=0}^k p_i^{-1} p_i \right) z^{-1} = 1.$$

## The Gray-Kambites construction (2)

RDG, Kambites (JEMS, to appear): When  $T = G$  (a group given by a finite special monoid pres.) and  $S = H$  (a f.g. subgroup), then

$$U(M_{G,H}) \cong H.$$



## The Gray-Kambites construction (3)

So, what is the RU-monoid of  $M_{T,S}$ ?

IgD, RDG (2024): RC-presented by  $p_i, q_i (= zp_i^{-1})$  ( $0 \leq i \leq k$ ),  $a^{(i)} (= p_i a p_i^{-1})$  ( $a \in A, 0 \leq i \leq k$ ),  $b^{(z)} (= z b z^{-1})$  ( $b \in B$ ), and relations

$$q_i w^{(i)} p_i = q_0 w^{(0)} p_0 \quad (w \in A^*, i = 1, \dots, k)$$

$$q_i u^{(i)} = q_i v^{(i)} \quad (u, v \in A^* \text{ s.t. } u = v \text{ holds in } T, \\ i = 0, 1, \dots, k)$$

$$q_i b^{(i)} = b^{(z)} q_i \quad (b \in B, i = 0, 1, \dots, k)$$

**NB.** For all  $u, v \in B^*$  s.t.  $u = v$  holds in  $S$ ,  $u^{(z)} = v^{(z)}$  can be RC-derived. In fact,  $\langle b^{(z)} : b \in B \rangle \cong S$ .

## The Gray-Kambites construction (4)

For example, when we take  $T = \{a\}^*$  and  $S = \langle \emptyset \rangle = \{1\}$  (and a silly presentation for  $T$ , say  $a = a$ , to have  $k = 1$ ) we get the RU-monoid

$$\text{MonRC}\langle a_1, a_1, p_0, p_1, q_0, q_1 \mid q_0 a_0^n p_0 = q_1 a_1^n p_1 \ (n \geq 0) \rangle.$$

This can be shown to be:

- ▶ **not** finitely RC-presented,
- ▶ with a **trivial** group of units.

**Conclusion:** There are non-finitely RC-presented RU-monoids out there!

Thank you! 😊❤️

---